

# Tropical method for proving plethysm

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Today's talk is based on my papers

- ① *Tropical integrable systems and Young tableaux: shape equivalence and Littlewood-Richardson correspondence*, Journal of Integrable Systems **3** (2018), no. 1, xyy011.
- ② *Jeu de taquin, uniqueness of rectification and ultradiscrete KP*, Journal of Integrable Systems **4** (2019), no. 1, xyz012.

# Young tableaux (Young 盤)

## Young tableaux of shape $\lambda$ (形 $\lambda$ の Young 盤)

$$\begin{array}{|c|c|} \hline 1 & 2 \\ \hline 3 & 3 \\ \hline \end{array}, \quad \lambda = (2, 2).$$

$$\begin{array}{|c|c|c|c|} \hline 1 & 2 & 2 & 4 \\ \hline 2 & 3 & 4 & \\ \hline 3 & & & \\ \hline 4 & & & \\ \hline \end{array}, \quad \lambda = (4, 3, 1, 1).$$

- A Young tableau is obtained by putting a natural number in each box of a Young diagram  $\lambda = (\lambda_1 \geq \lambda_2 \geq \dots)$ .
- Weakly increasing from left to right.
- Strongly increasing from top to bottom.

# Skew Young tableaux (歪 Young 盤)

## Skew Young tableaux of shape $\lambda/\mu$

$$\begin{array}{|c|c|} \hline & 1 & 2 \\ \hline 3 & 3 & \\ \hline \end{array}, \quad \lambda = (3, 2), \mu = (1)$$

$$\begin{array}{|c|c|} \hline & 1 & 4 \\ \hline 2 & 2 & \\ \hline 1 & & \\ \hline 3 & & \\ \hline \end{array}, \quad \lambda = (4, 3, 1, 1), \mu = (2, 1).$$

- Weakly increasing from left to right.
- Strongly increasing from top to bottom.

# Tropical method

	Field	Tropical semi-field
Language	Language of fields	Language of semi-fields
Addition	$a + b$	$\min[a, b]$
Multiplication	$a \cdot b$	$a + b$
Division	$a/b$	$a - b$

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Addition	$a + b$	$\min[a, b]$
Multiplication	$a \cdot b$	$a + b$
Division	$a/b$	$a - b$
Polynomial	$\sum_i a_i x^i$	$\min_i [a_i + i \cdot x]$
Rational function	$\frac{\sum_i a_i x^i}{\sum_j b_j x^j}$	$\min_i [a_i + i \cdot x]$ $- \min_j [b_j + j \cdot x]$

# Tropicalization

## Tropicalization

{Subtraction-free rational func.}  $\rightarrow$  {Tropical rational func.}

$$f = \frac{\sum_i a_i x^i}{\sum_j b_j x^j} \mapsto \hat{f} = \min_i [a_i + i \cdot x] - \min_j [b_j + j \cdot x]$$

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- 1  $\widehat{f + g} = \min[\widehat{f}, \widehat{g}]$ ,
- 2  $\widehat{f \cdot g} = \widehat{f} + \widehat{g}$ ,
- 3 If  $f$  and  $g$  are subtraction-free, the composition  $f \circ g$  is also subtraction-free:

$$\widehat{f \circ g} = \widehat{f} \circ \widehat{g}.$$



# Combinatorics of Young tableaux

Today we consider the following two combinatorial algorithms:

- ① Row insertion ( = Row bumping, "行挿入")
- ② Jeu de taquin (A french word means "15 パズル")

# Combinatorics of Young tableaux

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- The row insertion is characterized by a tropical recursion formula called **Tropical tableaux** (Bernstein-Kirillov [1] 2001, Noumi-Yamada [7] 2004), which is a tropical analogue of **discrete Toda equation**.
- Jeu de taquin is characterized by the **tropical KP equation** (Mikami [6] 2006, Katayama-Kakei [5] 2015), which is a tropical analogue of **discrete KP equation**.

# Row insertion

## Example

1	3	4	5	← 3
2	4	6		
4				

- The procedure starts from the top row.

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- The procedure starts from the top row.
- Insert 3 to the rightmost “possible position” as far as it does not violate the rule of Young tableau.

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- The procedure starts from the top row.
- Insert 3 to the rightmost “possible position” as far as it does not violate the rule of Young tableau.
- The “bumped” number 4 proceeds to the next row.



## Tropical tableaux (Cf. [1],[7])

The row insertion to a one row tableau is characterized by the tropical recursion formula

$$\begin{cases} y_i + b_i = a_i + x_i \\ b_{i+1} = a_{i+1} + X_{i+1} - X_i \\ X_i := \min_{r=1}^i [\sum_{k=1}^{r-1} (x_{i-k+1} - a_{i-k})] \end{cases}$$

where  $x_i$  (*resp.*  $y_i$ ) is the number  $i$ 's in the tableau before (*resp.* after) the insertion, and

$$a_i = \begin{cases} 1 & (\text{if } i \text{ is inserted}) \\ 0 & (\text{otherwise}) \end{cases} \quad b_i = \begin{cases} 1 & (\text{if } i \text{ is bumped}) \\ 0 & (\text{otherwise}) \end{cases}$$

# Tropical tableaux

## Example

1	3	6	7
2	4	8	
4			

← 3

row insertion  
 $\Rightarrow$

1	3	3	7
2	4	8	
4			

← 6

# Tropical tableaux

## Example

$$\begin{array}{|c|c|c|c|} \hline 1 & 3 & 6 & 7 \\ \hline 2 & 4 & 8 & \\ \hline 4 & & & \\ \hline \end{array} \leftarrow 3 \quad \text{row insertion} \Rightarrow \quad \begin{array}{|c|c|c|c|} \hline 1 & 3 & 3 & 7 \\ \hline 2 & 4 & 8 & \\ \hline 4 & & & \\ \hline \end{array} \leftarrow 6$$

$$\begin{aligned} (x_i)_i &= (1, 0, 1, 0, 0, 1, 1, 0, \dots), & (y_i)_i &= (1, 0, 2, 0, 0, 0, 1, 0, \dots), \\ (a_i)_i &= (0, 0, 1, 0, 0, 0, 0, 0, \dots), & (b_i)_i &= (0, 0, 0, 0, 0, 1, 0, 0, \dots), \end{aligned}$$

$$\begin{cases} a_i + x_i = y_i + b_i, & (\text{"\#}\{i\}\text{'s} \text{ remains unchanged"}) \\ b_{i+1} = a_i + X_{i+1} - X_i, \\ X_i = \min_{r=1}^i [\sum_{k=1}^{r-1} (x_{i-k+1} - a_{i-k})]. \end{cases}$$

# Tropical tableaux

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$$X_1 = \min[0] = 0,$$

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$$X_1 = \min[0] = 0, \quad X_2 = \min[0, x_2 - a_1] = 0,$$

# Tropical tableaux

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$$\begin{aligned}X_1 &= \min[0] = 0, & X_2 &= \min[0, x_2 - a_1] = 0, \\X_3 &= \min[0, x_3 - a_2, x_3 + x_2 - a_2 - a_1] = 0,\end{aligned}$$

# Tropical tableaux

$$\begin{aligned}(x_i)_i &= (1, 0, 1, 0, 0, 1, 1, 0, \dots), & (y_i)_i &= (1, 0, 2, 0, 0, 0, 1, 0, \dots), \\(a_i)_i &= (0, 0, 1, 0, 0, 0, 0, 0, \dots), & (b_i)_i &= (0, 0, 0, 0, 0, 1, 0, 0, \dots),\end{aligned}$$

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# Tropical tableaux

$$\begin{aligned}(x_i)_i &= (1, 0, 1, 0, 0, 1, 1, 0, \dots), & (y_i)_i &= (1, 0, 2, 0, 0, 0, 1, 0, \dots), \\(a_i)_i &= (0, 0, 1, 0, 0, 0, 0, 0, \dots), & (b_i)_i &= (0, 0, 0, 0, 0, 1, 0, 0, \dots),\end{aligned}$$

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$$(X_i)_i = (0, 0, 0, -1, -1, 0, 0, 0, \dots)$$

The equation  $b_{i+1} = a_i + X_{i+1} - X_i$  matches with the combinatorial description of the row insertion.



# Ultradiscrete Toda equation

Noumi-Yamada [7] showed that the row insertion algorithm is obtained from the **discrete Toda equation**;  $(x_i, a_i) \mapsto (y_i, b_i)$

$$\begin{cases} y_i b_i = a_i x_i \\ y_i + b_{i+1} = a_i + x_{i+1} & (i \geq 2) \\ b_2 = a_1 + x_2 \end{cases}$$

by **ultradiscretization** (= tropicalization).

# Ultradiscrete Toda equation

$$\begin{cases} y_i b_i = a_i x_i \\ y_i + b_{i+1} = a_i + x_{i+1} \quad (i \geq 2) \\ b_2 = a_1 + x_2 \end{cases}$$

$$b_2 = a_1 + x_2 = a_1 \cdot \left(1 + \frac{x_2}{a_1}\right)$$

$$b_3 = a_2 + x_3 - y_2 = a_2 + x_3 - \frac{a_2 x_2}{a_1 + x_2} = a_2 \cdot \frac{1 + \frac{x_3}{a_2} + \frac{x_3 x_2}{a_2 a_1}}{1 + \frac{x_2}{a_1}}$$

$$b_4 = a_3 + x_4 - y_3 = \dots = a_3 \cdot \frac{1 + \frac{x_4}{a_3} + \frac{x_4 x_3}{a_3 a_2} + \frac{x_4 x_3 x_2}{a_3 a_2 a_1}}{1 + \frac{x_3}{a_2} + \frac{x_3 x_2}{a_2 a_1}}$$

# Ultradiscrete Toda equation

Discrete Toda equation is rewritten as

$$\begin{cases} y_i b_i = a_i x_i \\ b_{i+1} = a_i \cdot \frac{\sum_{r=1}^{i+1} \prod_{k=1}^{r-1} \frac{x_{i-k+2}}{a_{i-k+1}}}{\sum_{r=1}^i \prod_{k=1}^{r-1} \frac{x_{i-k+1}}{a_{i-k}}} \end{cases}$$

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Tropicalizing this, we have

$$\begin{cases} a_i + x_i = y_i + b_i, \\ b_{i+1} = a_i + X_{i+1} - X_i, \\ X_i = \min_{r=1}^i [\sum_{k=1}^{r-1} (x_{i-k+1} - a_{i-k})]. \end{cases}$$

# Row insertion as a tropical map

	Field	Tropical
$x_i$	input at a discrete time $t$	$\#(i\text{'s in a row})$
$a_i$	"	$\#(i\text{'s to be inserted})$
$y_i$	output at a discrete time $t$	$\#(i\text{'s in a row after insertion})$
$b_i$	"	$\#(i\text{'s proceed to the next row})$
	$\begin{cases} a_i x_i = y_i b_i \\ a_i + x_{i+1} = y_i + b_{i+1} \\ b_2 = a_1 + x_2 \end{cases}$ (discrete Toda equation)	$\begin{cases} a_i + x_i = y_i + b_i \\ b_{i+1} = a_i + X_{i+1} - X_i \end{cases}$ (Row insertion)

(cf. Noumi-Yamada [7])

## Example

		1	2
		2	3
	1	3	4
2	4	4	

- Choose an “inside corner.”

# Jeu de taquin

## Example

		1	2
		2	3
	1	3	4
2	4	4	

- Choose an “inside corner.”

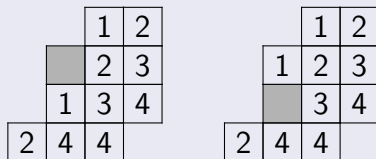
## Example

		1	2
		2	3
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- Choose an “inside corner.”
- Compare numbers on the right and below, and slide the smaller one. (If they are equal, move the below one).



## Example



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		1	2
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		1	2
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# Jeu de taquin

## Example

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		3	4
2	4	4	

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	1	2	3
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2	4	4	

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# Jeu de taquin as a tropical map

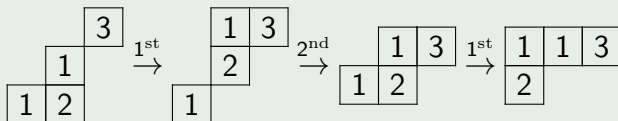
	Field	Tropical
$f_{i,j}^t$	input of the discrete KP eq.	$\# \left( \begin{array}{l} 0, 1, 2, \dots, j \text{'s contained} \\ \text{in the } 1^{\text{st}}, 2^{\text{nd}}, \dots, i^{\text{th}} \text{ rows} \\ \text{after the } t^{\text{th}} \text{ move} \end{array} \right)$
	$f_{i,j}^t f_{i,j+1}^{t+1}$ $= f_{i+1,j+1}^t f_{i-1,j}^{t+1} + f_{i,j+1}^t f_{i,j}^{t+1}$ (discrete KP equation)	$F_{i,j}^t + F_{i,j+1}^{t+1}$ $= \max[F_{i+1,j+1}^t + F_{i-1,j}^{t+1},$ $F_{i,j+1}^t + F_{i,j}^{t+1}]$ (Jeu de taquin)

(cf. Katayama-Kakei [5])

# Rectification

A **rectification** of a skew tableau is a (non-skew) Young tableau that is obtained by a sequence of Jeu de taquin.

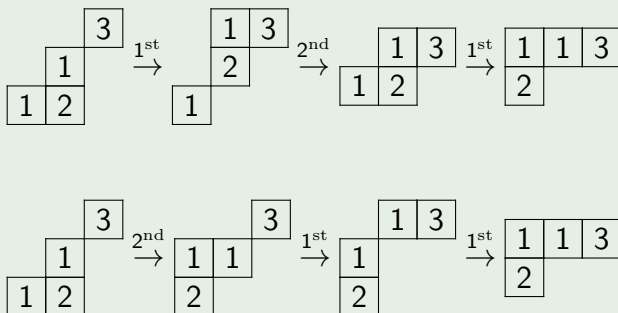
## Example



# Rectification

A **rectification** of a skew tableau is a (non-skew) Young tableau that is obtained by a sequence of Jeu de taquin.

## Example



## Theorem (Uniqueness of rectification, cf.[2])

A rectification of a skew tableau does not depend on the choice of a sequence of jeu de taquin.

Combinatorial proofs are well-known (cf. Fulton's textbook "Young Tableau [2]")

## Theorem (Uniqueness of rectification, cf.[2])

A rectification of a skew tableau does not depend on the choice of a sequence of jeu de taquin.

Combinatorial proofs are well-known (cf. Fulton's textbook "Young Tableau [2]")

Today I give (an outline of) an alternative proof of this theorem by using [tropical integrable systems](#).



Language of fields

$$p: A \rightarrow B \quad (\text{Subtraction-free rational map})$$

(Geometrization)  $\uparrow$   $\downarrow$  (Tropicalization)

Language of semi-fields

$$\hat{p}: \hat{A} \rightarrow \hat{B}$$

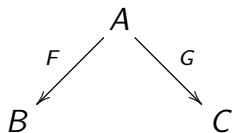
(A piecewise linear map obtained from  $p$  by  $\left\{ \begin{array}{l} + \mapsto \min \\ \times \mapsto + \end{array} \right.$ )

# Main trick

Language of fields

$\uparrow$     $\downarrow$

Language of semi-fields (=combinatorial objects)



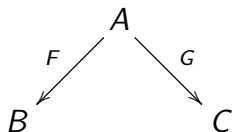
$F, G$ : Maps that are defined combinatorially  
(e.g. row insertion, jeu de taquin)

# Main trick

Language of fields

(Lift)  $\uparrow$   $\downarrow$

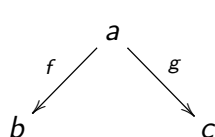
Language of semi-fields (=combinatorial objects)



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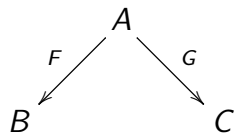
Language of fields



$$\hat{f} = F, \hat{g} = G$$

(Lift)  $\uparrow$   $\downarrow$

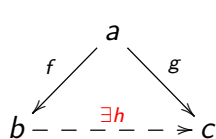
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# Main trick

Language of fields

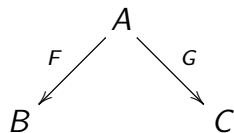


$$\widehat{f} = F, \widehat{g} = G$$

$h$ : Subtraction-free

(Lift)  $\uparrow$   $\downarrow$

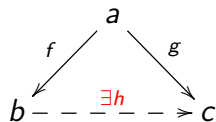
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$F, G$ : Maps that are defined combinatorially  
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# Main trick

Language of fields

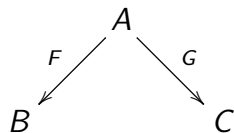


$$\widehat{f} = F, \widehat{g} = G$$

$h$ : Subtraction-free

(Lift)  $\uparrow$        $\downarrow$  (Tropicalize)

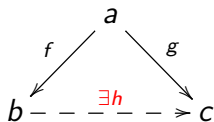
Language of semi-fields (=combinatorial objects)



$F, G$ : Maps that are defined combinatorially  
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# Main trick

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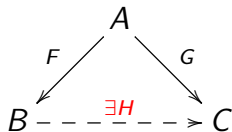


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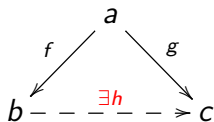
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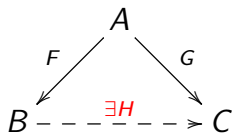


$$\widehat{f} = F, \widehat{g} = G$$

$h$ : Subtraction-free

(Lift)  $\uparrow$        $\downarrow$  (Tropicalize)

Language of semi-fields (=combinatorial objects)

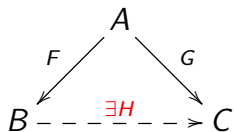


$F, G$ : Maps that are defined combinatorially  
(e.g. row insertion, jeu de taquin)

$H$ : Well-defined map



# Main trick

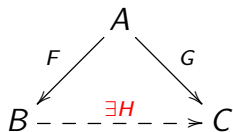


$F, G$ : Maps that are well-known

$H$ : A map just defined

I am happy enough because:

# Main trick



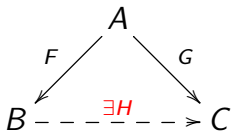
$F, G$ : Maps that are well-known

$H$ : A map just defined

I am happy enough because:

“There exists a map  $H$  that makes the diagram commutative.”

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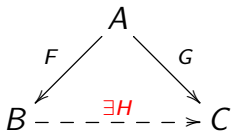
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= “The procedure of taking an inverse image of  $F$  and sending it by  $G$  does not depend on the choice of the inverse image. ”

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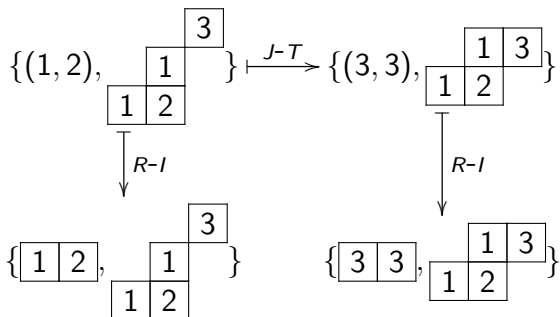
“Is the procedure  $X$  independent of the choice of  $Y$  ?” is a FAQ for those who study the combinatorics of Young tableaux.

# Outline of proof

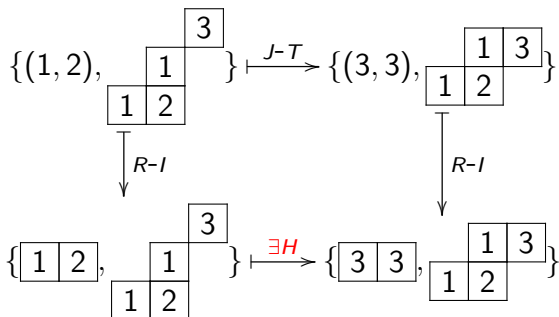
Previous researches have shown that the two maps  $J-T$  (Noumi-Yamada 2004) and  $R-I$  (Katayama-Keiei 2015) are given by a tropical rational map.

$$\left\{ (1, 2), \begin{array}{|c|c|} \hline & 3 \\ \hline 1 & 2 \\ \hline \end{array} \right\} \xrightarrow{J-T} \left\{ (3, 3), \begin{array}{|c|c|} \hline 1 & 3 \\ \hline 1 & 2 \\ \hline \end{array} \right\}$$
$$\left\{ (1, 2), \begin{array}{|c|c|} \hline & 3 \\ \hline 1 & 2 \\ \hline \end{array} \right\} \xrightarrow{R-I} \left\{ \begin{array}{|c|c|} \hline 1 & 2 \\ \hline \end{array}, \begin{array}{|c|c|} \hline & 3 \\ \hline 1 & 2 \\ \hline \end{array} \right\}$$

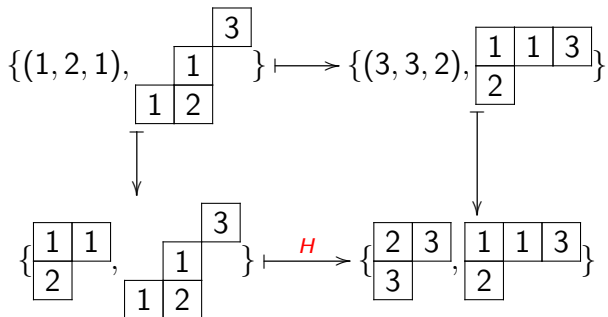
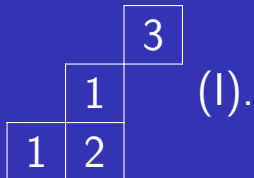
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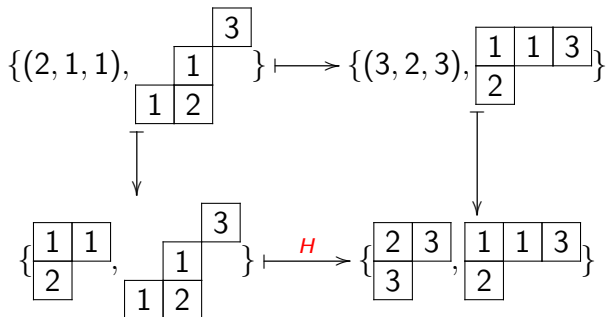
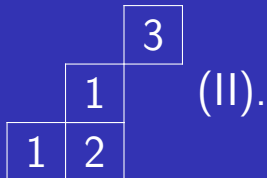


# Two rectifications of





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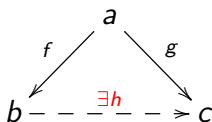


# Uniqueness

In each case, we have

$$H \left( \left( \begin{array}{|c|c|} \hline 1 & 1 \\ \hline 2 & \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline & & 3 \\ \hline & 1 & \\ \hline 1 & 2 & \\ \hline \end{array} \right) = \left( \begin{array}{|c|c|} \hline 2 & 3 \\ \hline 3 & \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline 1 & 1 & 3 \\ \hline 2 & & \\ \hline \end{array} \right).$$

This is an essential reason why a rectification is unique.



The subtraction-free map  $h$  is constructed by using the [factorization of totally positive matrices](#), which is a popular technique in the theory of discrete integrable system (cf. Iwao [3, 4])

# References I

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- [5] Yosuke Katayama and Saburo Kakei, *Jeu de taquin slide and ultradiscrete KP equation (in Japanese)*, Reports of RIAM Symposium **26AO-S2** (2015), 133–138.
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- [7] Masatoshi Noumi and Yasuhiko Yamada, *Tropical Robinson-Schensted-Knuth correspondence and birational Weyl group actions*, Representation theory of algebraic groups and quantum groups (T. Shoji, M. Kashiwara, N. Kawanaka, G. Lusztig, and K. Shinoda, eds.), vol. 40, Soc. Japan, Tokyo, 2004, pp. 371–442.